

CS 520: Introduction to Artificial Intelligence

Prof. Louis Steinberg

Lecture 8:

Knowledge-based Agents

Propositional Logic

Review - Lisp

Symbol value vs meaning as function

- Value:

- Assign with `(setf var <expression>)`
- Create local binding with `let` or `let*`

```
(let* ((var <expr>)
      (var <expr>) ... )
  <expr> ... )
```

- Meaning as function

- Set global meaning with `defun`
- Create local scope with `flet` or `labels`

Review

Functions as arguments

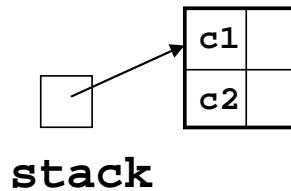
- *Value* of `#'foo` is *function meaning* of `foo`
- `(funcall fn-expr argx1 ... argxn)`
means treat *value* of `fn-expr` as a function, call it with values of `argx1 ... argxn` as arguments

Review - Lexical Scope

- **Scope of a variable is *syntactic* scope of `let`, `defun`, etc.**
- **If a function is created, it carries bindings with it**
 - **Closure: represents function, points to binding & code**

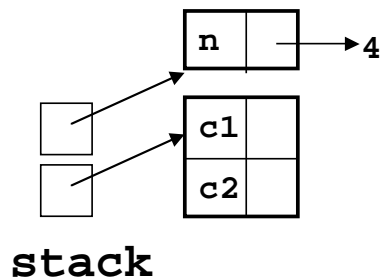
Closures

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(defun counter (n)
  #'(lambda (i)
      (setf n(+ n i))))
(let((c1 (counter 4))
     (c2 (counter 0)))
  (list
   (funcall c1 10)
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   (funcall c1 60)))
=> (14 3 5 74)
```



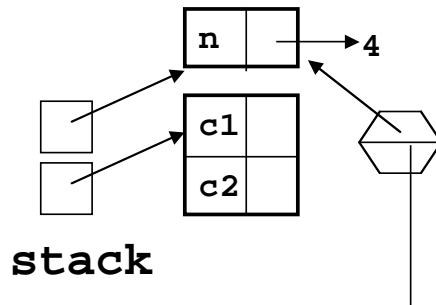
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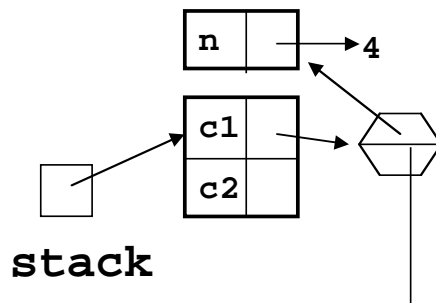
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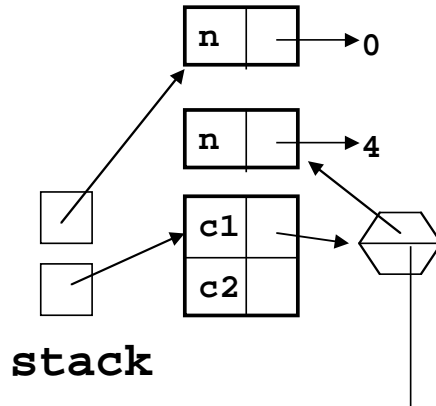


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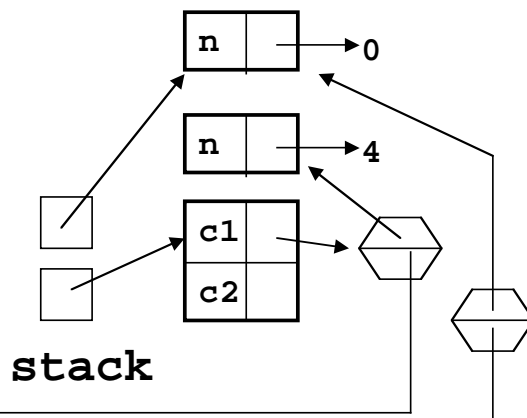


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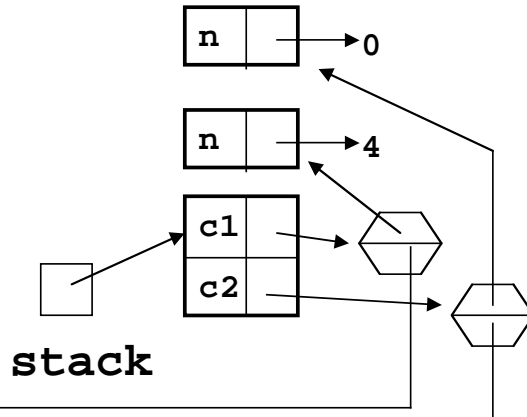
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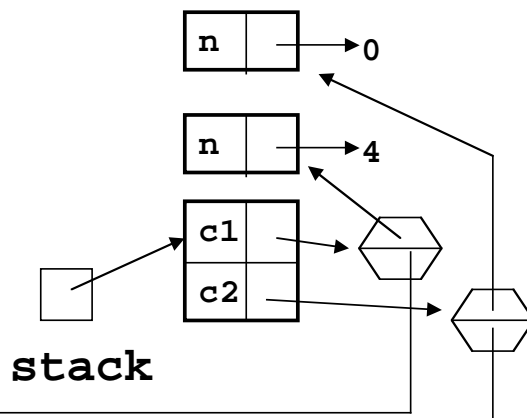
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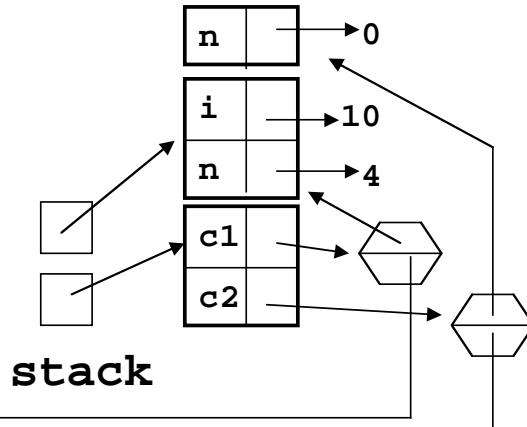
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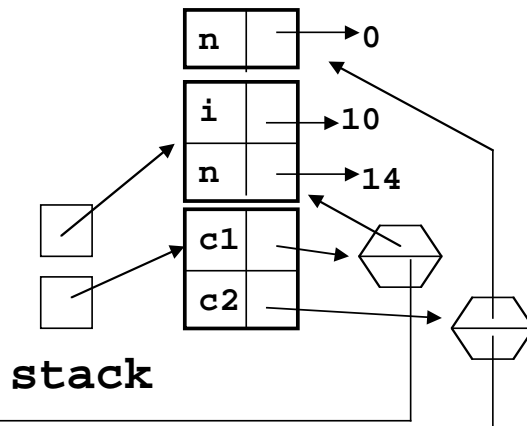
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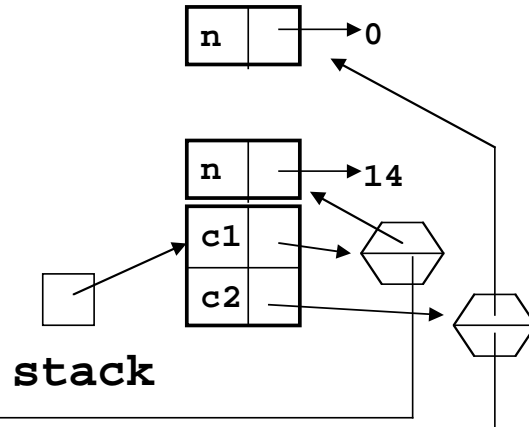
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Programming Techniques

- **Data-driven programming**
- **To write a program for X, extend Lisp to be a good language for X**

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Lecture 08

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Quiz

- **Write code for some.**
(some fn list) calls fn on each element of list until a call returns non-nil, then returns that value.
If all calls return nil, some returns nil.
(some #'oddp '(2 4 5)) => t
(some #'oddp '(2 4 6))=> nil

Mechanisms for Intelligence

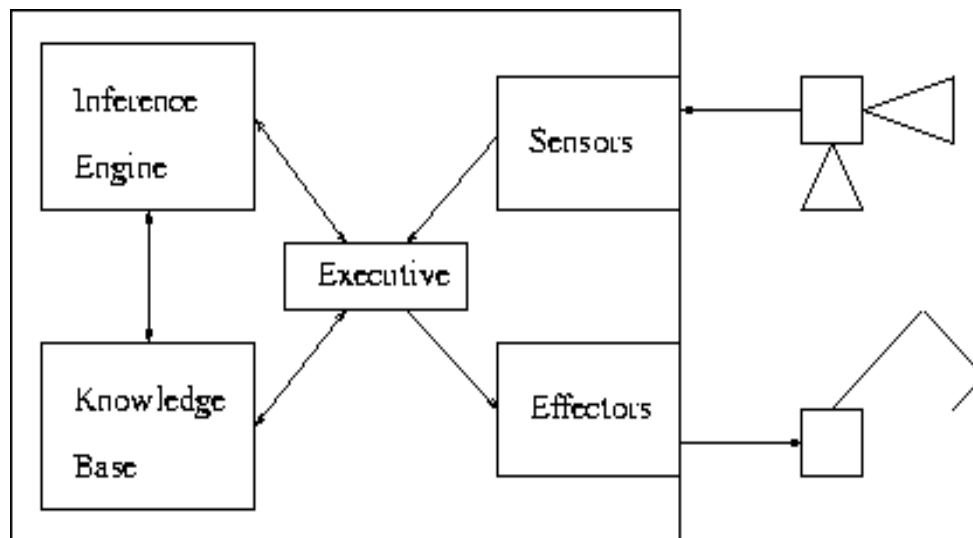
Given a problem, formalize it as an instance of a general class of problems, which you know how to solve

- **Search**
 - Water jugs => graph search
 - N Queens => constraint satisfaction
 - Chess => minimax tree
- **Inference**

Knowledge-Based Agents

- A knowledge-based agent has knowledge of the world (its environment) and knowledge of its actions.
- Using logical reasoning, it maintains a description of the world, and will try to determine an action that will meet its goals.

Knowledge-based Agent Architecture



Assertions and Queries

- **Executive asserts** (tells KB) “I sense”
- **Executive queries** (asks KB) “What should I do?”
- **Internal assertions and queries of KB**
 - Sense and Do
 - State of the world
- **Eg:** (sense obstacle +45 degrees)
(sense obstacle -45 degrees)
(sense no obstacle 0 degrees)
=> (facing door)
=> (do move-forward)

Why Inference?

- **Ideal:**
 - Tell agent *what* is true
 - Fact1, fact2, ...
 - Give agent a goal
 - Agent decides for itself *how* to use the facts to attain the goal
- **Better than giving agent a program**
 - More autonomous
 - For given effort we get agent with more general competence.

Logic

- **A logic consists of**
 - **Formal language: sentences**
 - Represent statements about the world
 - **Inference mechanism:**
 - $\{\text{sentence}_1, \dots, \text{sentence}_n\} \Rightarrow \text{sentence}_{n+1}$
 - Creates new sentences
- **Eg, propositional logic**

Propositional Logic

- **Language:**
 - Symbols: a, b, \dots
 - Operators: $\neg, \wedge, \vee, \Rightarrow, ()$
 - **Semantics:**
 - Symbol represents a statement about world that is either true or false
 - Operators combine truth values
- $a = \text{“it is raining”}, b = \text{“the sidewalk is wet”}$
 $a \wedge b = \text{“it is raining and the sidewalk is wet”}$

Models

- **Model:**
 - An assignment of real-world meaning to the symbols of a logic
- A sentence in a logic may be true or false with respect to a model
- A sentence may be such that it is *true* in every model: “tautology”
- A sentence may be such that it is *true* in some model: “satisfiable”
- A sentence may be such that it is *false* in every model: “contradiction”

Validity of Inference

- An inference rule
 $\{\text{old sentences}\} \rightarrow \text{new-sentence}$
is *valid* if for every model it is true that
 - If $\text{meaning}(\text{old-sentences})$ is true
then $\text{meaning}(\text{new-sentence})$ must be true
- IE, $(\text{old-sentences} \Rightarrow \text{new-sentence})$ is a tautology
 - Old-sentences *entail* new sentence

Soundness and Completeness

- An inference mechanism is **sound** or **truth-preserving** if it generates from the KB only sentences that are entailed by the KB.
- **Tracing through the steps of deriving a sentence from a KB using a sound inference mechanism is called a proof .**
- **If the inference mechanism can derive any sentence that is entailed by the KB, the mechanism is said to be complete .**

Propositional Inference

- **We can test entailment in propositional logic with a truth table**

$(A \wedge B)$ entails $\neg(\neg A \vee \neg B)$

A	B	$(A \wedge B)$	$\neg(\neg A \vee \neg B)$
T	T	T	T
T	F	F	F
F	T	F	F
F	F	F	F

Inference by Rules

- **Problem:** truth table size is exponential in number of variables
- **Solution:** test entailment by searching for a sequence of inference rule applications

Rules of Inference for Propositional Logic

Modus Ponens : $\alpha; \alpha \Rightarrow \beta / \beta$

And-elimination : $\alpha_1 \wedge \alpha_2 \wedge \dots \wedge \alpha_n / \alpha_1$

And-introduction ...

Or-elimination ...

Or-introduction ...

Double-negation-elimination

Unit resolution

Resolution

...

Resolution

- **Problem: too many rules -> high branching factor for search**
- **Want single complete rule**
 - **Resolution**