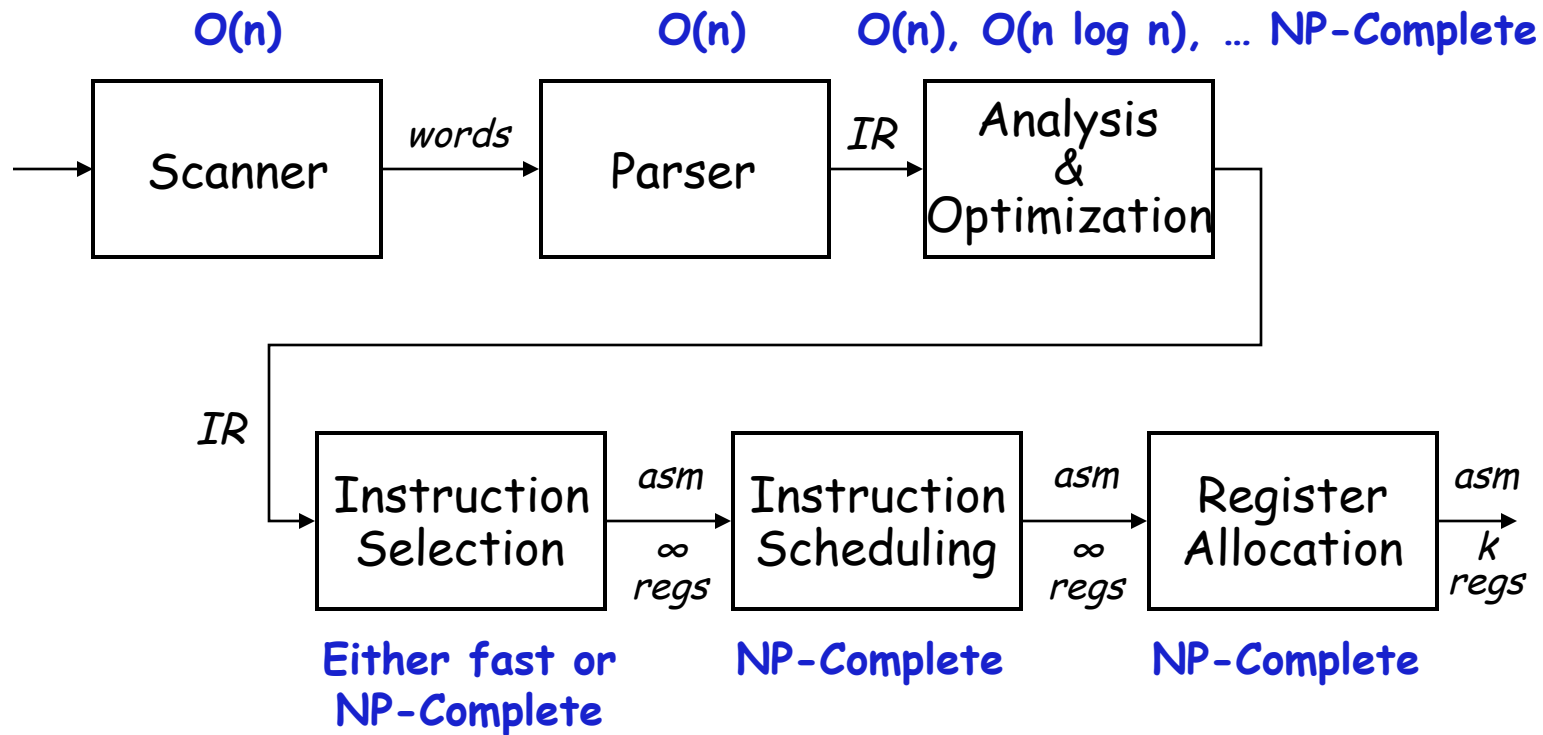




CS415 Compilers

Code Generation

These slides are based on slides copyrighted by
Keith Cooper, Ken Kennedy & Linda Torczon at Rice
University



- A compiler is a lot of fast stuff followed by some hard problems
- The hard stuff is mostly in **code generation** and **optimization**
 - For superscalars, its allocation & scheduling that is particularly important

Conventional wisdom says that we lose little by solving these problems independently

Instruction selection

- Use some form of pattern matching
- Assume enough registers or target “important” values

Note: many fuzzy terms here!

Instruction scheduling

- Within a block, list scheduling is “close” to optimal
- Across blocks, build framework to apply list scheduling

Optimal for
> 85% of blocks

Register allocation

- Start from virtual registers & map into k
- Focus on good priority heuristic

Definition

- All those nebulous properties of the code that impact performance & code “quality”
- Includes code, approach for different constructs, cost, storage requirements & mapping, & choice of operations
- Code shape is the end product of many decisions *(big & small)*

Impact

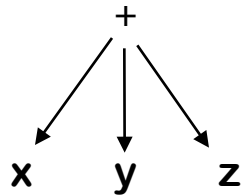
- Code shape influences algorithm choice & results
- Code shape can encode important facts, or hide them

Rule of thumb: expose as much derived information as possible

- Example: explicit branch targets in ILOC simplify analysis

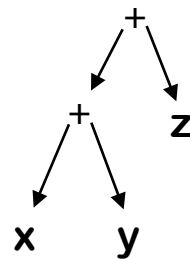
An example

$$x + y + z$$



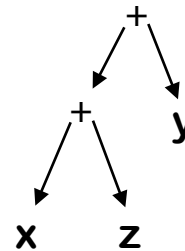
$$x + y \rightarrow t1$$

$$t1 + z \rightarrow t2$$



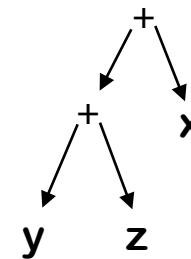
$$x + z \rightarrow t1$$

$$t1 + y \rightarrow t2$$



$$y + z \rightarrow t1$$

$$t1 + x \rightarrow t2$$



- What if x is 2 and z is 3?
- What if $y+z$ is evaluated earlier?

Addition is commutative & associative for integers

The “best” shape for $x+y+z$ depends on contextual knowledge
 → There may be several conflicting options

Another example -- the case statement

- Implement it as cascaded if-then-else statements
 - Cost depends on where your case actually occurs
 - $O(\text{number of cases})$
- Implement it as a binary search
 - Uniform ($\log n$) cost
- Implement it as a jump table
 - Lookup address in a table & jump to it
 - Uniform (constant) cost

Compiler must choose best implementation strategy

No amount of massaging or transforming will convert one into another

The key code quality issue is holding values in registers

- When can a value be safely allocated to a register?
 - When only 1 name can reference its value (**no aliasing**)
 - Pointers, parameters, aggregates & arrays all cause trouble
- When should a value be allocated to a register?
 - When it is both safe & profitable

Encoding this knowledge into the *IR* (*register-register model*)

- Use code shape to make it known to every later phase
- Assign a virtual register to anything that can go into one
- Load or store the others at each reference

Relies on a strong register allocator

More code generation

Read EaC: Chapter 7.1 - 7.5