

## TEST 1 - Some Solutions

1. The challenge of this question is to find a known theorem, or fact, which supports the assertion. If you cannot, you may believe the assertion is not generally true, and then search for a counterexample. Constructing a counterexample for the assertion that is still consistent with the given data may actually take some time. If you don't see an immediate reason why the assertion may be false, a good strategy might be to say "I believe this is not generally true, and will come back later to construct a counter-example".

(a)  $P(Y = 3) < 1$  is **FALSE**: the constant random variable  $Y = 3$  is consistent with the given data.

(b)  $P(X \geq 1) \geq \frac{1}{10}$  is **TRUE** by the analog of Markov's inequality. It says that  $Z$  is a random variable with  $E(Z) = \mu$ ,  $Z \leq B$  with probability 1, and  $t \leq \mu$ , then

$$P(Z \geq t) \geq \frac{\mu - t}{B - t}$$

Here we have  $E(X) = 2$ ,  $t = 1$  and  $B = 10$ . The right hand side is  $1/9 \geq 1/10$ .

(c)  $V(X+Y) = 9 + V(Y)$  is **TRUE** by calculation:  $V(X+Y) = E[(X+Y)^2] - (E(X+Y))^2$  and this is  $V(X) + V(Y) + 2[E(XY) - E(X)E(Y)]$ .

(d)  $P(X \geq 8) \leq 1/4$  is **FALSE**. The temptation is to invoke Markov's inequality, but we are NOT GIVEN the necessary fact that  $X$  is a non-negative random variable. If  $X$  is 4 or  $-1$ , each with probability  $1/2$ , the required data for  $X$  are satisfied but not the stated probability inequality.

(e)  $P(X \geq 5) \leq 2/5$  is **TRUE** by Tchebycheff's inequality

(f) **FALSE** -  $X$  and  $Y$  are NOT necessarily independent, just because  $E(XY) = E(X)E(Y)$ . I will try to put in a simple counterexample later.

2. (a)

$$2 \frac{(k!)(k!)}{(2k)!}$$

(b)

$$\frac{\binom{2k}{k} k!}{(2k)!}$$

There are  $\binom{2k}{k}$  ways to determine the locations of the odds. Given the locations, there is only one ordering for the odds - they have to increase. There are  $k!$  orderings of the evens in the remaining locations.

For the algorithm, generate  $(\pi_1, \dots, \pi_{2k})$  a random permutation. Initialize another array  $A$  of size  $2k$  to zero. The first  $k$  entries in  $\pi$  point to where the evens go, so set  $A[\pi_j] \leftarrow 2j$ ,  $j = 1, \dots, k$ . The last  $k$  entries of the permutation point to where the odds go. So just scan the  $A$  array and set  $2j - 1$  in place of the  $j^{\text{th}}$  zero in  $A$ ,  $j = 1, \dots, k$ . The alg is linear.

(c)

$$\frac{((k-1)!)^2}{(2k)!}$$

(d)

$$\frac{\binom{2k}{k}(k-1)!}{(2k)!}$$

$P_D(C) = 0 \neq P(C)$ , so they are NOT independent.

3. I gave 5 pts here only if some info was given

4. Randomly sample.

5. (a)  $G$  is a cycle and a “tail”. The min-cut is the tail separating the vertex of degree one from the other  $n - 1$  vertices.

(b)  $2/n$ .