Internet Technology

10. Multicast Routing

Paul Krzyzanowski

Rutgers University

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Broadcast & Multicast

• Broadcast routing
  – Deliver a packet to all nodes in the network

• Multicast routing
  – Deliver a packet to some subset of nodes in the network
N-Way Unicast

• The initial sender (source node) makes $N$ copies of a datagram, one for each destination node, and transmits them
  – Potentially a lot of overhead: $N$ copies go over the first link
    • It would be more efficient for a router to make copies for each link
  – How does it know all the recipients?
Uncontrolled Flooding

• A router will make a copy of the packet and forward to all neighbors
• Cycles will result in a broadcast storm
  – Endless replication of packets
Controlled Flooding

• Sequence number controlled flooding
  – Sender places its source address and a broadcast sequence number into the packet
  – Each node keeps a list of \{source address, sequence number\} of each packet that was forwarded
  – Before copying & forwarding a packet, check the list
    • If we saw it, drop it

• Reverse path forwarding (RPF)
  – Packet is duplicated & forwarded ONLY IF it was received via the link that is the shortest path to the sender
  – Shortest path is found by checking the forwarding table to the source address
Reverse Path Forwarding (RPF)

Drop packets that do not originate from the link for the shortest path back to the source.

Example:
A message is initially sent to router A. It duplicates it and sends it to routers B & C.

Router B duplicates it & sends it to its outgoing links: to C, D, and E.

Each router checks whether the link that the message arrived on is part of the shortest path to the destination.

When C receives the message from B, it checks that the shortest path for the message is through A, not C. Hence, the message is rejected.
Ensure that every node receives only one copy of the broadcast packet.

Create an overlay network that is a subset of the connected network:
- contains all nodes in the network
- subset of edges
- connected graph
- no cycles

This is a spanning tree

If each link has a cost
\[
\text{cost(tree)} = \text{sum of link costs}
\]
tree whose cost is the minimum of all possible spanning trees is a minimum spanning tree

A source node would send a packet only onto the links that are part of the spanning tree

A node only has to keep track of which of its neighbors are part of the spanning tree
Building a spanning tree

• Center-based approach

• Define a center node (rendezvous point)

• Nodes send tree-join messages to this center node
  – The message is forwarded toward the center until it
    • arrives at a node that already belongs to the spanning tree
    • or arrives at the center
  – The path that the tree-join message traversed defines a branch of
    the spanning tree
Example: building a spanning tree

Arbitrarily pick E as the core.

F sends the first tree-join message

EF link becomes a branch
Arbitrarily pick E as the core.

F sends the first tree-join message
EF link becomes a branch

B sends the next tree-join message
BD, DE become branches
(could have been BC, CE)
 Arbitrarily pick E as the core.

 F sends the first tree-join message  
  \( EF \) link becomes a branch

 B sends the next tree-join message  
  \( BD, DE \) become branches  
  (could have been \( BC, CE \))

 A sends the next tree-join message  
 This one goes through C.  
  \( AC \) and \( CE \) become branches
Arbitrarily pick E as the core.

F sends the first tree-join message  
*EF* link becomes a branch

B sends the next tree-join message  
*BD, DE* become branches  
(could have been *BC, CE*)

A sends the next tree-join message  
This one goes through C.  
*AC and CE* become branches

If a node connected to C needs to join,  
C is already a part of the tree.
Arbitrarily pick E as the core.

F sends the first tree-join message
*EF* link becomes a branch

B sends the next tree-join message
*BD, DE* become branches
(could have been *BC, CE*)

A sends the next tree-join message
This one goes through C.
*AC* and *CE* become branches

G sends the next (last) tree-join message
*GD* becomes a link (no need to build the tree to *E*, since *D* is already part of it)
IP multicast routing
Multicast routing

• Deliver messages to a subset of nodes
• How do we identify the recipients?
  – Enumerate them in the header?
    • What if we don’t know?
    • What if we have thousands of recipients?
• Use a special address to identify a group of receivers
  – A copy of the packet is delivered to all receivers associated with that group
  – Class D multicast IP address
    • 32-bit address that starts with 1110
      \(224.0.0.0/4 = 224.0.0.0 – 239.255.255.255\)
  – Host group = set of machines listening to a particular multicast address
IP multicasting

• Can span multiple physical networks

• Dynamic membership
  – Machine can join or leave at any time

• No restriction on number of hosts in a group

• Machine does not need to be a member to send messages

• Efficient: Packets are replicated only when necessary
IP multicast addresses

• Addresses chosen arbitrarily for an application

• Well-known addresses assigned by IANA
  – Internet Assigned Numbers Authority
  – See
    http://www.iana.org/assignments/multicast-addresses/multicast-addresses.xml
  – Similar to ports – service-based allocation
    • For ports, we have:
      – FTP: port 21, SMTP: port 25, HTTP: port 80
    • For multicast, we have:
      224.0.0.1: all systems on this subnet
      224.0.0.2: all multicast routers on subnet
      224.0.23.173: Philips Health
      224.0.23.52: Amex Market Data
      224.0.12.0-63: Microsoft & MSNBC
IGMP

- Internet Group Management Protocol (IGMP)
  - Operates between a host and its attached router
  - Goal: allow a router to determine to which of its networks to forward IP multicast traffic
  - IP protocol (protocol number 2)

- Three message types
  - Membership_query
    - Sent by a router to all hosts on an interface to determine the set of all multicast groups that have been joined by the hosts on that interface
  - Membership_report
    - Host response to a query or an initial join or a group
  - Leave_group
    - Host indicates that it is no longer interested
    - Optional: router infers this if the host does not respond to a query
Multicast Forwarding

• IGMP allows a host to *subscribe* to a multicast stream
• What about the source?
  – There is no protocol for the source!
  – It just sends to a class D address
  – Routers have to do the work
IGMP & Wide-Area Multicast Routing

Internet multicast routing

IGMP

PIM

no protocol!
Multicast Forwarding

• **IGMP**: Internet Group Management Protocol
  – Designed for routers to talk with hosts on directly connected networks

• **PIM**: Protocol Independent Multicast
  – Multicast Routing Protocol for delivering packets across routers
  – Topology discovery is handled by other protocols
Flooding: Dense Mode Multicast

- Relay multicast packet to all connected routers
  - Use a spanning tree and use reverse path forwarding (RPF) to avoid loops
  - Feedback & cut-off if there are no receivers on a link
    - A router sends a prune message.
    - Periodically, routers send messages to refresh the prune state
  - Flooding is initiated by the sender’s router

- Advantage:
  - Simple
  - Good if the packet is desired in most locations

- Disadvantage:
  - wasteful on the network, wasteful extra state & packet duplication on routers
Sparse Mode Multicast

• Initiated by the routers at each receiver

• Each router needs to ask for a multicast feed with a PIM Join message
  – Initiated by a router at the destination that gets an IGMP join
  – Spanning tree constructed
    • Join messages propagate to a defined rendezvous point
    • Sender transmits only to the rendezvous point
  – A Prune message stops a feed

• Advantage
  – Packets go only where needed
  – Creates extra state in routers only where needed
IP Multicast in use

• Initially exciting:
  – Internet radio, NASA shuttle missions, collaborative gaming

• But:
  – Few ISPs enabled it
  – For the user, required tapping into existing streams
    (not good for on-demand content)
  – Industry embraced unicast instead
IP Multicast in use: IPTV

• IPTV is emerging as the biggest user of IP multicast

• Cable TV systems: aggregate bandwidth ~ 4.5 Gbps
  – Video streams: MPEG-2 or MPEG-4 (H.264)
  – MPEG-2 HD: ~30 Mbps
  – MPEG-4 HD: ~6-9 Mbps; DVD quality: ~2 Mbps
IP Multicast in use: IPTV

• Traffic is within the provider’s network
  – QoS: typically mix of ATM and/or IP
    • 2.5 Mbps VBR video
    • 256 kbps CBR voice
    • Remainder: ABR for IP traffic
  – Unicast for video on demand
  – Multicast for live content
    • When you select a channel, you join a multicast group via IGMPv2
      – Local office checks if you are authorized.
      – If yes, routers add the user to the group
    • Burst of unicast data to get the I-frame to ensure 150 msec channel switching times.
The end