

CS344: Design And Analysis of Computer Algorithms

S. Muthu Muthukrishnan

Details

- **Meeting:** Tues, Fri. 11.30 – 12.50 PM. SEC 118
- **Professor:** Muthu Muthukrishnan, 445 2379,
Core 319.
- **Office hours:** Tuesdays 3.30 – 4.30 PM.
- **Contact:**
 - muthu at cs.rutgers.edu
 - <http://www.cs.rutgers.edu/~muthu>
- **TAs:**

More Details

- **Book:**
 - Cormen, Leiserson, Rivest, Stein.
<http://theory.lcs.mit.edu/~clr/>
- **Grading:**
 - 2 Midterms 25% each: March 4, April 8.
 - 1 Final Exam 40%: May 6.
 - Homeworks 10%.
 - Teams of X people each.
 - Assignment given on Fridays, due the following Friday.
- **Policy:** Writing “**I don’t know**” gives you 25% of the points.
- Khachiyan is an excellent lecturer! Both sections are independent.

Syllabus

- Asymptotic analysis of algorithms
- Recursion, induction, recurrence equations.
- Sorting and Selection.
- Hashing and balanced search structures
- Graphs and Graph algorithms
- Dynamic Programming
- String matching.
- Polynomials, matrices.
- NP Completeness
- Assorted topics.

Big Picture of Computer Science

- **Theory**
 - What is computable? Turing machines, recursive functions.
 - What is efficiently computable? Algorithms, Complexity theory.
 - Abstract math, practical methods.
- **Systems**
 - Databases, OS, Networks, Programming Languages, Web,...
 - Theory, tools.
- **Popular Science**
 - CPU MIPS, GUIs, Blogs, SPAM,...

- Design of efficient algorithms for problems.
- Mathematically analyze running time and space used.
- Lower bounds: What is the best possible algorithm, what is “best”?
- Are there solvable problems that have no efficient algorithms?

Example

- Problem: Maximum subarray sum
- Input: Array $A[1..n]$ of integers.
- Output: $[i,j]$ such that $A[i] + \dots + A[j]$ is largest.

Naïve Strategy

- For all pairs $1 \leq i \leq j \leq n$, compute
 - $\text{sum}(i,j) = A[i] + \dots + A[j]$
- Simple triple loop will do.
- Running time:
 - Roughly n^3 .

```
For i=1 to n
  For j=i+1 to n
    Sum[i,j]=0;
    For k=i to j
      Sum[i,j]+=A[k].
```

Less Naive Strategy

- When you have $\text{sum}(i,j)=A[i]+\dots+A[j]$, it is easy to get $\text{sum}(i,j+1)=A[i]+\dots+A[j+1]$.

```
For i=1..n
  Sum[i,i]=A[i]
  For j=i+1 .. n
    Sum[i,j]=Sum[i,j-1]+A[j]
```

- Double loop will do.
- Running time: $n+(n-1)+\dots+1$, roughly n^2 .

What is the fastest running time possible?